# Lecture 8

Parallel Programming

#### PLQ Review: traits and enums

```
trait Shape {
fn area(\&self) -> f32;
}
impl Shape for Circle {
fn area(\&self) -> f32 {
    self.0 * self.0 * 3.14
 }
}
impl Shape for Rect {
fn area(\deltaself) -> f32 {
   self.0 * self.1 }
}
```

```
enum Shape {
 Circle(Circle),
Rect(Rect),
}
impl Shape {
 fn area(\&self) -> f32 {
    match self {
      &Shape::Circle(Circle(r)) =>
       r * r * 3.14,
     \&Shape::Rect(Rect(x, y)) =>
       x * y }
 }
}
```
### PLQ Review: traits and enums

```
trait Shape {
 fn area(\text{dself}) -> f32;
}
```
Can extend implementers but not functionality

- implement trait for a new type
- can't add new trait method without breaking existing implementers

```
enum Shape {
Circle(Circle),
Rect(Rect),
}
```
Can extend functionality but not implementers

- Add new enum method to add functionality
- Can't add new variant to enum without breaking existing methods

# Where are we? Where are we going?

Multi-threading is the motivation behind many of Rust's features. Today, we'll start tying together many features:

- limiting mutability
- smart pointers
- ownership
- traits

Coming up:

- Two lectures on parallelism
- One lecture on **unsafe Rust**
- One lecture on Rust ecosystem
- One flex lecture

# Who cares about parallelism?

Easy to buy more cores, impossible to buy faster cores

● How to use more cores?



# Why is parallelism hard?

```
#define NUM_THREADS 10
#define INCREMENTS 10000
int counter = 0;
void* increment counter(void* arg) {
   for (int i = 0; i < INCREMENTS; i++) {
      counter = counter + 1;
 }
   return NULL;
```

```
int main() {
   pthread t threads[NUM_THREADS];
   for (int i = 0; i < NUM THREADS; i++) {
       pthread create(&threads[i], NULL,
                      increment counter, NULL);
 }
  // wait for threads to finish...
```

```
 printf("Counter: %d\n", counter);
```
What is printed?

}

}

#### Data race

Multiple accesses with at least one writer

```
for (int i = 0; i < INCREMENTS; i++) {
 int cur count = counter;
 int new count = cur count + 1;
counter = new count;
}
```
7

Thread 1:

Thread 2:



time

# Mutex ("mutual exclusion")

```
#define NUM_THREADS 10
#define INCREMENTS 10000
```
int counter =  $0$ ;

pthread mutex t counter mutex;

```
void* increment counter(void* arg) {
   for (int i = 0; i < INCREMENTS; i++) {
      pthread mutex lock(&counter mutex);
       counter++;
       pthread mutex unlock(&counter mutex);
 }
    return NULL;
}
              "critical section"
```

```
int main() {
   pthread t threads[NUM_THREADS];
   pthread mutex init(&counter mutex, NULL);
   for (int i = 0; i < NUM THREADS; i++) {
       pthread create(&threads[i], NULL,
                      increment counter, NULL);
 }
    // wait for threads to finish...
```

```
 printf("Counter: %d\n", counter);
```
}

# Critical section

Guarantee that only one thread will be executing critical section at a time

● Other threads wait if mutex is locked

```
for (int i = 0; i < INCREMENTS; i++) {
lock (mutex) ;
 int cur count = counter;
 int new count = cur count + 1;
 counter = new count;
unlock (mutex)
}
```
Thread 1:

Thread 2:



# Mutex challenges

What are bugs you can make when using mutexes?

# Mutex challenges

- Forget to use one at all
- Forget to lock
- Forget to unlock
- Lock in wrong order
- Lock while already locked
- Unlock while already unlocked

# Rust's bold claims

Impossible to forget to protect data with a mutex

● Compile-time quarantee of no data races!

Impossible to

- Forget to lock
- Forget to unlock
- Unlock while already unlocked

Can still do:

- Lock in wrong order
- Lock while already locked

Initially [safety and concurrency] seemed orthogonal, but to our amazement, the solution turned out to be identical: **the same tools that make Rust safe also help you tackle concurrency head-on**."

Compiler enforces rules for safe concurrency. **"Thread safety isn't just documentation; it's law."**

*https://blog.rust-lang.org/2015/04/10/Fearless-Concurrency.html*

#### What does parallelism look like in Rust?

}

use std::thread;

const NUM THREADS: usize = 10; const INCREMENTS PER THREAD: usize = 10000;

```
fn main() {
  let mut counter = 0;
  for in 0..NUM_THREADS {
   thread::spawn(|| {
           for _ in 0..INCREMENTS_PER_THREAD {
              counter += 1:
 }
       });
 }
    // wait for threads to finish...
   println!("Counter: {}", counter);
```
#### What does parallelism look like in Rust?

}

use std::thread;

const NUM THREADS: usize = 10; const INCREMENTS PER THREAD: usize = 10000;

*Same data race problem as before!*

Which (if any) of Rust's ownership/borrowing rules are violated?

```
fn main() {
   let mut counter = 0;
   for __ in 0..NUM_THREADS {
       thread::spawn(|| {
           for _ in 0..INCREMENTS PER THREAD
               counter += 1; }
        });
\qquad \qquad \}wait for threads to finish...
                                                important 
                                                part
                                                thread executes 
                                                a closure
```
println!("Counter: {}", counter);

#### One (of many) problems

```
error[E0499]: cannot borrow `counter` as mutable more than once at a time 
 --> race.rs:10:36 
 | 
10 | thread::spawn(|| { 
                           | ^^ `counter` mutably borrowed in previous iteration of loop 
11 | for in 0..INCREMENTS PER THREAD {
12 | counter += 1; 
                      | ------- borrows occur due to use of `counter` in closure 
13 | } 
14 | });
```
Need to be able to mutate counter, but can't give out multiple mutable references?

# One (of many) problems

```
error[E0373]: closure may outlive `main`, but borrows `counter`, which is owned by `main` 
 | 
10 | thread::spawn(|| { 
                          | ^^ may outlive borrowed value `counter` 
11 | for in 0..INCREMENTS PER THREAD {
12 | counter += 1; 
                   | ------- `counter` is borrowed here 
 | 
 note: function requires argument type to outlive `'static`
```
How long can a thread live? Who owns counter?

#### Questions to solve:

Need to be able to mutate counter, but can't give out multiple mutable references?

How long can a thread live? Who owns counter?

#### One (of many) problems

```
help: to force the closure to take ownership of `counter`, use the `move` keyword<br>I
10 | let handle = thread::spawn(move || {
error[E0373]: closure may outlive `main`, but borrows `counter`, which is owned by `main` 
 | 
10 | thread::spawn(|| { 
                       | ^^ may outlive borrowed value `counter` 
11 | for in 0..INCREMENTS PER THREAD {
12 | counter += 1; 
                 | ------- `counter` is borrowed here 
 | 
 note: function requires argument type to outlive `'static` 
  --> race.rs:10:22 
 | 
10 | let handle = thread::spawn(|| { 
 | ______________________^ 
11 | | for in 0..INCREMENTS PER THREAD
12 | | counter += 1; 
13 | | } 
14 | | }); 
 | |__________^ 
 | 
 | ++++
```
#### Rc/RefCell to the rescue?

}

use std::rc::Rc; use std::cell::RefCell; use std::thread;

```
const NUM THREADS: usize = 10;
const INCREMENTS PER THREAD: usize = 10000;
```
Shared ownership and interior mutability

```
fn main() {
   let counter = Rc: new(RefCell: new(0)) ;
   for in 0..NUM THREADS {
       let counter = Rc:: clone(&counter) ;
        thread::spawn(move || {
           for in 0..INCREMENTS PER THREAD {
               *counter.borrow mut() += 1;
 }
        });
 }
    // wait for threads to finish...
   println!("Counter: {}", counter<mark>.borrow mut()</mark>);
```
#### Rc/RefCell to the rescue?



#### Rc/RefCell to the rescue?

#### use std::rc::Rc; **error cannot be std::rc**ell::Reference between the same safely safely safely safely safely safely safely safely safely

 $l_{\rm c}$ : new ( $l_{\rm c}$ : n

 **|**  $\left[ -\right] % \includegraphics[width=0.9\textwidth]{images/TrDiM-Architecture.png} \caption{The first two different values of $d=2$ and $d=1/2$ (left) and the second two different values of $d=1/2$ (right) and the second two different values of $d=1/2$ (right).} \label{TrDiS}$ let counter counters and containing the counter-14 **and 14 control to the committee of the stands for instruments** 

The type  $Rc \le T$  provides shared ownership of a value of type T, allocated in the heap. Invoking cl ter to the same allocation in the heap. When the last  $\mathop{\rm RC}\nolimits$  pointer to a given allocation is destroyed, the value stored in  $\mathop{\rm tr}\nolimits$  $\rm{ed.}$  . The  $\sim$  15  $\pm$  15  $\pm$  15  $\pm$  16  $\pm$ construction provides shared entitled per a ratio **EXECUTE:** *PCCTS provides shared ownership of a value of type T allow*  **| | |** allocation (often referred to as "inner value") is also dropped.

ed references in Rust disallow mutation by default, and Rc is no except ething inside an Rc. If you need mutability, put a Cell or RefCell **17 | |** }

 $f_n$  mainly  $\left(\lambda\right)$  (

Bo uses non-atomic reference counting. This means that overhe wait for the same for the finished to finish the finish of the finish of the finish of the finish of the finish between threads. If you need multi-threaded, atomic reference counting, use sync:: Arc. be shown about the complete counting. This means that overhead is very long bat an no cannot be s **|**

}

# Shared ownership and interior mutability

**Single Threaded:**

**Multi-Threaded:**

Rc: shared ownership with a reference count

Arc: shared ownership with an *atomic* reference count

RefCell: interior mutability by panicking if multiple mutable borrows

Mutex: interior mutability by making other threads wait

#### Rust parallelism third attempt

}

```
use std::sync::{Arc,Mutex};
use std::thread;
```

```
const NUM THREADS: usize = 10;
const INCREMENTS PER THREAD: usize = 10000;
```

```
Mutex provides interior mutability!
Exclusive ownership while mutex locked
```

```
fn main() {
   let counter = \text{Arc::new}(\text{Mutes::new}(0)) ;
   for __ in 0..NUM_THREADS {
        let counter = \text{Arc::}clone(&counter);
        thread::spawn(move || {
            for in 0..INCREMENTS PER THREAD {
                 let mut guard = \frac{\text{counter}, \text{lock}}{\text{!}} . unwrap();
                 *quard += 1; }
         });
 }
    // wait for threads to finish...
   println!("Counter: {}", counter.lock().unwrap());
```
# Zooming in on thread::spawn



(1) make shared resource

(2) clone shared resource handle once per thread

(3) move shared resource handle into thread

Very common pattern when spawning threads

# Shared-state concurrency

Threads have separate stacks but a shared heap

Share ownership of *same* mutex

Control access of shared counter via mutex



# Zooming in on Mutexes

```
use std::sync::{Mutex, MutexGuard};
fn increment(m: &Mutex<u32>) {
  let guard: MutexGuardXu32 = m.lock().unwrap();
   *guard += 1; // guard is dropped -> mutex is unlocked
}
fn main() {
```

```
let counter = Mutes: new(0); increment(&counter);
```
}

Why is it impossible to forget to unlock a mutex?

Why is it impossible to forget to unlock a mutex?

```
fn copy_inner(m: &Mutex<u32>) -> u32 {
  let guard: MutexGuard<u32> = m.lock().unwrap();
  let value = *guard;
   return value
   // guard is dropped -> mutex is unlocked
}
```
Impossible to not drop MutexGuard

Why is it impossible to forget to lock a mutex?

Why is it impossible to forget to lock a mutex?

Only lock gives access to inner value

```
fn no_lock1(m: &Mutex<u32>) -> u32 {
   let value = \frac{m}{2}??();
    return v;
}
```
Can't have reference to inner value outlive guard

```
fn no_lock2(m: &Mutex<u32>) -> &mut u32 {
   let mut guard = m. lock().unwrap();
   let v: \text{kmut } u32 = \text{kmut } \star \text{quard}; return v;
```
}

Why is it impossible to unlock a mutex twice?

Why is it impossible to unlock a mutex twice? fn double\_unlock (m: &Mutex<u32>) -> u32 {

```
let guard = m. lock().unwrap();
  let value = *quard;
  drop(guard);
  drop (guard) ;
   return value;
}
```
# Quiz code

https://godbolt.org/z/P179e9a3z

#### Compared to C



pthread mutex t counter mutex;



let counter: Mutex<u32> = Mutex::new(0);

Nothing associates mutex with data!

Mutex has ownership of protected data

# Quiz takeaways

Ownership and reference lifetimes make it impossible to misuse a mutex

> Initially [safety and concurrency] seemed orthogonal, but to our amazement, the solution turned out to be identical: **the same tools that make Rust safe also help you tackle concurrency head-on**."

Compiler enforces rules for safe concurrency. **"Thread safety isn't just documentation; it's law."**

*https://blog.rust-lang.org/2015/04/10/Fearless-Concurrency.html*

#### How does the compiler know?



suspicious: is this analysis special-cased for standard library types?

#### How does the compiler know?



# Send and Sync

Send: it is safe to send this type to another thread

Sync: it is safe to share this type between threads

● T is Sync if and only if &T is Send

"safe" -> won't cause memory safety errors or data races

#### Types that are…



#### Not Sync or Send: Rc

use std::rc::Rc; use std::thread;

```
fn main() {
```

```
let count = Rc::new(0); let clone1 = Rc::clone(&count);
 let clone2 = Rc::clone(&count);
 thread::spawn(move || {
     drop(clone1);
 });
```

```
 thread::spawn(move || {
     drop(clone2);
```
});

}



#### Not Sync or Send: Rc



# Send but not Sync: RefCell

```
use std::{thread, cell::RefCell};
fn increment(r: &RefCell<u32>) {
   let mut count = r.borrow mut();
   *count += 1;}
```

```
fn main() {
   let count = RefCell::new(11);let ref1 = \&count;let ref2 = \&count: thread::spawn(move || {
        increment(ref2);
    });
    increment(ref1);
```
}



borrows is non-atomic: two threads could both successfully borrow\_mut at the same time

#### Send but not Sync: RefCell



# Send and Sync are special

Automatically derived for all types whose members are Send/Sync

You won't implement Send/Sync for your types, but you may use them as bounds for type parameters in generic functions

If the lock is always automatically released, is it possible to have a deadlock in Rust?

If the lock is always automatically released, is it possible to have a deadlock in Rust?

Yes! Double lock, as shown before

What else?

If the lock is always automatically released, is it possible to have a deadlock in Rust? Yes!

```
fn swap1(a: Arc<Mutex<u32>>, b: Arc<Mutex<u32>>) {
   let mut guard_a = a.\text{lock}().\text{unwrap}();
   let mut guard b = b.\text{lock}().\text{unwrap}();
    // do the swap
}
fn swap2(a: Arc<Mutex<u32>>, b: Arc<Mutex<u32>>) {
   let mut guard_b = b.\,lock().\unumber{unwrap()};
   let mut guard a = a.\text{lock}().\text{unwrap}();
    // do the swap
}
                                               \leftarrow thread 1 waiting
                                              thread 2 waiting
```
If the lock is always automatically released, is it possible to have a deadlock in Rust?

```
fn swap1(a: Arc<Mutex<u32>>, b:
Arc<Mutex<u32>>) {
   let mut guard a = a.\text{lock}().\text{unwrap}();
   let mut guard b = b.\text{lock}().\text{unwrap}();
    // do the swap
}
```

```
fn swap2(a: Arc<Mutex<u32>>, b:
Arc<Mutex<u32>>) {
   let mut guard b = b.\text{lock}().\text{unwrap}();
   let mut guard a = a.\text{lock}().\text{unwrap}();
    // do the swap
}
```
#### Yes!

}

```
fn main() {
   let a = Arc::new(Mutes::new(10));let b = Arc::new(Mutes::new(20));
```

```
let a cloned = Arc::clone(\&a);let b cloned = Arc::clone(\&b); thread::spawn(move || {
    swap1(a cloned, b cloned);
 });
swap2(Arc::clone(&a), Arc::clone(&b));
```
#### One last form of interior mutability

use std::sync::atomic::{AtomicUsize, Ordering};

```
fn increment_atomic(counter: &AtomicUsize) {
  counter.fetch_add(1, Ordering::SeqCst);
```
}

**Atomics**: allow mutation through a shared reference.

Other threads are guaranteed not to observe intermediate values.

What's this?

● it's complicated: just use Ordering:SeqCst