Correlated Equilibria

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- But maybe there is some richer family of equilibria we can shoot for...
- Analogous to our earlier relaxation from Pure to Mixed equilibria.

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But one player never gets any utility...

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- 3. So both players play STOP with probability p = 100/101, and play GO with probability (1 p) = 1/101.
- 4. This is even worse! Now both players get payoff 0 in expectation (rather than just one of them), and risk a horrific negative utility.

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Worse: there is no set of mixed strategies that creates this distribution over action profiles.



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- 5. We can generalize this...

Correlated Equilibrium

Definition

A correlated equilibrium is a distribution \mathcal{D} over action profiles A such that for every player i, every action \hat{a}_i , and every action a_i^* :

$$\mathrm{E}_{a \sim \mathcal{D}}[u_i(a)|a_i = \hat{a}_i] \geq \mathrm{E}_{a \sim \mathcal{D}}[u_i(a_i^*, a_{-i})|a_i = \hat{a}_i]$$

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For example: Conditioned on seeing STOP, you know your opponent will GO, so STOP is a best response. Conditioned on seeing GO, you know your opponent will STOP, so GO is a best response.

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A coarse correlated equilibrium is a distribution \mathcal{D} over action profiles A such that for every player i, and every action a_i^* :

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- 4. The difference: the suggestion just has to be a best response on average, not *conditioned* on having seen it.
- Whether it is sensible depends on whether you have to commit to following the correlating device up front, or have the option of deviating after seeing the suggestion.



CCE can occasionally suggest obviously bad actions. CE cannot.

Consider:

		Α		A B		C			
A		(1,1)		(1,1) (-1)		(-1,-1)		(0,0)	
В	((-1,-1)		$) \mid (1,1)$		(0,0)			
C		(0,0)		(0,0) (-1.1,-1.1		,-1.1)		
			1	4	В	С			
		Α	1,	/3					
		В	,		1/3				
		C				1/3			

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C		(0,0)		(0,0)		(-1.1,-1.1)		
			Α		В	С		
		Α	1,	/3				
		В			1/3			
		C				1/3		

The payoff for each player for playing according to this distribution is:

$$(1/3) \cdot 1 + (1/3) \cdot 1 - (1/3) \cdot 1.1 = 0.3$$

		А		В		С			
Α		(1,1)		(1,1)		(-1,-1)		(0,0)	
В	((-1,-1)		(1,1)		(0,0)			
C		(0,0)		(0,0)	(-1.1,	-1.1)		
			1	4	В	С			
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the payoff a player would get by playing the fixed action A or B while his opponent randomized would be:

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and the payoff for playing C would be strictly less than zero.

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the payoff a player would get by playing the fixed action A or B while his opponent randomized would be:

$$(1/3) \cdot 1 - (1/3) \cdot 1 + (1/3) \cdot 0 = 0$$

and the payoff for playing C would be strictly less than zero.

Hence this is a CCE *even though* conditioned on being told to play *C*, it is not a best response. This means that the given distribution is a coarse correlated equilibrium, *but not* a CE.



Solution Concept Recap

 $\textit{DSE} \subset \textit{PSNE} \subset \textit{MSNE} \subset \textit{CE} \subset \textit{CCE}$

Hierarchies

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1. Starting at MSNE, we have guaranteed existence.

Hierarchies

Solution Concept Recap

$$DSE \subset PSNE \subset MSNE \subset CE \subset CCE$$

- 1. Starting at MSNE, we have guaranteed existence.
- 2. Want to show: Starting at CE, we have computational tractability.

Definition

For a *strategy modification rule* $F_i: A_i \rightarrow A_i$ and an action profile $a \in A$:

$$\operatorname{Regret}_i(a, F_i) = u_i(F_i(a_i), a_{-i}) - u_i(a)$$

i.e. it is how much player i regrets not applying F_i to change his action.

We say that F_i is a constant strategy modification rule if $F_i(a_i) = F_i(a_i')$ for all $a_i, a_i' \in A_i$.

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We can give an alternative characterization of CCE:

Definition

A distribution \mathcal{D} is a coarse correlated equilibrium if for every player i and for every constant strategy modification rule F_i :

$$\mathbb{E}_{a \sim \mathcal{D}}[\operatorname{Regret}_i(a, F_i)] \leq 0$$

1. An immediate consequence of this definition is that if a^1,\ldots,a^T are a sequence of actions with $\Delta(T)$ regret, then $\bar{a}=\frac{1}{T}\sum_{t=1}^T a^t$ forms a $\Delta(T)$ -approximate coarse correlated equilibrium.

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- 2. In particular, if everyone plays an (arbitrary) game with the PW algorithm, after T steps they will have generated a sequence of plays that corresponds to a $\Delta(T) = 2\sqrt{\log k/T}$ -approximate CCE
- 3. Can we approach computing CE in the same way? First step: characterize CE in terms of regret.

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To see this:

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- 4. Look for learning algorithms with stronger regret guarantees...



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$$\frac{1}{T}\sum_{t=1}^{T}\ell_{\boldsymbol{s}^t} \leq \frac{1}{T}\sum_{t=1}^{T}\ell_{F(\boldsymbol{s}^t)} + \Delta(T)$$

for all strategy modification rules F and for $\Delta(T) = o(1)$.

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- 3. "No Swap Regret"
- 4. We'll see how to do this! (Next lecture).

Thanks!

See you next class!